

AperTO - Archivio Istituzionale Open Access dell'Università di Torino

Higher Hamming weights for locally recoverable codes on algebraic curves

This is the author's manuscript

Original Citation:

Availability:

This version is available <http://hdl.handle.net/2318/1560283> since 2016-04-27T17:03:20Z

Published version:

DOI:10.1016/j.ffa.2016.03.004

Terms of use:

Open Access

Anyone can freely access the full text of works made available as "Open Access". Works made available under a Creative Commons license can be used according to the terms and conditions of said license. Use of all other works requires consent of the right holder (author or publisher) if not exempted from copyright protection by the applicable law.

(Article begins on next page)

HIGHER HAMMING WEIGHTS FOR LOCALLY RECOVERABLE CODES ON ALGEBRAIC CURVES

Edoardo Ballico ^{*}& Chiara Marcolla¹

ABSTRACT

We study locally recoverable codes on algebraic curves. In the first part of the manuscript, we provide a bound on the generalized Hamming weight of these codes. In the second part, we propose a new family of algebraic geometric LRC codes, which are LRC codes from the Norm-Trace curve. Finally, using some properties of Hermitian codes, we improve the bounds on the distance proposed in [1] of some Hermitian LRC codes.

1 INTRODUCTION

The v -th generalized Hamming weight $d_v(C)$ of a linear code C is the minimum support size of v -dimensional subcodes of C . The sequence $d_1(C), \dots, d_k(C)$ of generalized Hamming weights was introduced by Wei [37] to characterize the performance of a linear code on the wire-tap channel of type II. Later, the GHWs of linear codes have been used in many other applications regarding the communications, as for bounding the covering radius of linear codes [15], in network coding [26], in the context of list decoding [7, 9], and finally for secure secret sharing [18]. Moreover, in [2] the authors show in which way an arbitrary linear code gives rise to a secret sharing scheme, in [16, 17] the connection between the trellis or state complexity of a code and its GHWs is found and in [4] the author proves the equivalence to the dimension/length profile of a code and its generalized Hamming weight. For these reasons, the GHWs (and their *extended* version, the *relative* generalized Hamming weights [21, 19]) play a central role in coding theory. In particular, generalized and relative generalized Hamming weights are studied for Reed-Muller codes [10, 23] and for codes constructed by using an algebraic curve [6]

* The first author is partially supported by MIUR and GNSAGA of INDAM (Italy).

^{*} *Department of Mathematics, University of Trento, Italy.* Email address: ballico@science.unitn.it

¹ *Department of Mathematics, University of Turin, Italy.* Email address: chiara.marcolla@unito.it

as Goppa codes [24, 38], Hermitian codes [12, 25] and Castle codes [27].

In this paper, we provide a bound on the generalized Hamming weight of locally recoverable codes on the algebraic curves proposed in [1]. Moreover, we introduce a new family of algebraic geometric LRC codes and improve the bounds on the distance for some Hermitian LRC codes.

Locally recoverable codes were introduced in [8] and they have been significantly studied because of their applications in distributed and cloud storage systems [3, 13, 32, 34, 35]. We recall that a code $C \in (\mathbb{F}_q)^n$ has locality r if every symbol of a codeword c can be recovered from a subset of r other symbols of c .

In other words, we consider a finite field $K = \mathbb{F}_q$, where q is a power of a prime, and an $[n, k]$ code C over the field K , where $k = \log_q(|C|)$. For each $i \in \{1, \dots, n\}$ and each $a \in K$ set $C(i, a) = \{c \in C \mid c_i = a\}$. For each $I \subseteq \{1, \dots, n\}$ and each $S \subseteq C$ let S_I be the restriction of S to the coordinates in I .

Definition 1. Let C be an $[n, k]$ code over the field K , where $k = \log_q(|C|)$. Then C is said to have **all-symbol locality r** if for each $a \in \mathbb{F}_q$ and each $i \in \{1, \dots, n\}$ there is $I_i \subset \{1, \dots, n\} \setminus \{i\}$ with $|I_i| \leq r$, such that for $C_{I_i}(i, a) \cap C_{I_i}(i, a') = \emptyset$ for all $a \neq a'$. We use the notation (n, k, r) to refer to the parameters of this code.

Note that if we receive a codeword c correct except for an erasure at i , we can recover the codeword by looking at its coordinates in I_i . For this reason, I_i is called a *recovering set* for the symbol c_i .

Let C be an (n, k, r) code, then the distance of this code has to verify the bound proved in [28, 8] that is $d \leq n - k - \lceil k/r \rceil + 2$. The codes that achieve this bound with equality are called *optimal* LRC codes [32, 34, 35]. Note that when $r = k$, we obtain the Singleton bound, therefore optimal LRC codes with $r = k$ are MDS codes.

LAYOUT OF THE PAPER This paper is divided as follows. In Section 2 we recall the notions of algebraic geometric codes and the definition of algebraic geometric locally recoverable codes introduced in [1]. In Section 3 we provide a bound on the generalized Hamming weights of the latter codes. In Section 4 we propose a new family of algebraic geometric LRC codes, which are LRC codes from the Norm-Trace curve. Finally, in Section 5 we improve the bounds on the distance proposed in [1] for some Hermitian LRC codes, using some properties of the Hermitian codes.

2 PRELIMINARY NOTIONS

2.1 Algebraic geometric codes

Let $K = \mathbb{F}_q$ be a finite field, where q is a power of a prime. Let \mathcal{X} be a smooth projective absolutely irreducible nonsingular curve over K . We denote by $K(\mathcal{X})$ the rational func-

tions field on \mathcal{X} . Let D be a divisor on the curve \mathcal{X} . We recall that the *Riemann-Roch space* associated to D is a vector space $\mathcal{L}(D)$ over K defined as

$$\mathcal{L}(D) = \{f \in K(\mathcal{X}) \mid (f) + D \geq 0\} \cup \{0\}.$$

where we denote by (f) the divisor of f .

Assume that P_1, \dots, P_n are rational points on \mathcal{X} and D is a divisor such that $D = P_1 + \dots + P_n$. Let G be some other divisor such that $\text{supp}(D) \cap \text{supp}(G) = \emptyset$. Then we can define the algebraic geometric code as follows:

Definition 2. The **algebraic geometric code** (or AG code) $C(D, G)$ associated with the divisors D and G is defined as

$$C(D, G) = \{(f(P_1), \dots, f(P_n)) \mid f \in \mathcal{L}(G)\} \subset K^n.$$

The dual $C^\perp(D, G)$ of $C(D, G)$ is an algebraic geometric code.

In other words an algebraic geometric code is the image of the evaluation map $\text{Im}(ev_D) = C(D, G)$, where the *evaluation map* $ev_D : \mathcal{L}(G) \rightarrow K^n$ is given by

$$ev_D(f) = (f(P_1), \dots, f(P_n)) \in K^n.$$

Note that if $D = P_1 + \dots + P_n$ and we denote by $\mathcal{P} = \{P_1, \dots, P_n\}$ we can also indicate ev_D as $ev_{\mathcal{P}}$.

2.2 Algebraic geometric locally recoverable codes

In this section we consider the construction of algebraic geometric locally recoverable codes of [1].

Let \mathcal{X} and \mathcal{Y} be smooth projective absolutely irreducible curves over K . Let $g : \mathcal{X} \rightarrow \mathcal{Y}$ be a rational separable map of curves of degree $r + 1$. Since g is separable, then there exists a function $x \in K(\mathcal{X})$ such that $K(\mathcal{X}) = K(\mathcal{Y})(x)$ and that x satisfies the equation $x^{r+1} + b_r x^r + \dots + b_0 = 0$, where $b_i \in K(\mathcal{Y})$. The function x can be considered as a map $x : \mathcal{X} \rightarrow \mathbb{P}_K$. Let $h = \deg(x)$ be the degree of x .

We consider a subset $S = \{P_1, \dots, P_s\} \subset \mathcal{Y}(K)$ of \mathbb{F}_q -rational points of \mathcal{Y} , a divisor Q_∞ such that $\text{supp}(Q_\infty) \cap \text{supp}(S) = \emptyset$ and a positive divisor $D = tQ_\infty$. We denote by

$$\mathcal{A} = g^{-1}(S) = \{P_{ij}, \text{ where } i = 0, \dots, r, j = 1, \dots, s\} \subset \mathcal{X}(K),$$

where $g(P_{ij}) = P_i$ for all i, j and assume that b_i are functions in $\mathcal{L}(n_i Q_\infty)$ for some natural numbers n_i with $i = 1, \dots, r$.

Let $\{f_1, \dots, f_m\}$ be a basis of the Riemann-Roch space $\mathcal{L}(D)$. By the Riemann-Roch Theorem we have that $m \geq \deg(D) + 1 - g_{\mathcal{Y}}$, where $g_{\mathcal{Y}}$ is the genus of \mathcal{Y} .

From now on, we assume that $m = \deg(D) + 1 - g_{\mathcal{Y}}$, where $\deg(D) = t\ell$, and we consider the K -subspace V of $K(\mathcal{X})$ of dimension rm generated by

$$\mathcal{B} = \{f_j x^i, i = 0, \dots, r-1, j = 1, \dots, m\}.$$

We consider the evaluation map $\text{ev}_{\mathcal{A}} : V \rightarrow K^{(r+1)s}$. Then we have the following theorem.

Theorem 1. *The linear space $C(D, g) = \text{Span}_{K^{(r+1)s}} \langle \text{ev}_{\mathcal{A}}(\mathcal{B}) \rangle$ is an (n, k, r) algebraic geometric LRC code with parameters*

$$\begin{aligned} n &= (r+1)s \\ k &= rm \geq r(\text{tl} + 1 - g_y) \\ d &\geq n - \text{tl}(r+1) - (r-1)h. \end{aligned}$$

Proof. See Theorem 3.1 of [1]. □

The AG LRC codes have an additional property. They are LRC codes (n, k, r) with $(r+1) \mid n$ and $r \mid k$. The set $\{1, \dots, n\}$ can be divided into $n/(r+1)$ disjoint subsets U_j for $1 \leq j \leq s$ with the same cardinality $r+1$. For each i the set $I_i \subseteq \{1, \dots, n\} \setminus \{i\}$ is the complement of i in the element of the partition U_j containing j , i.e. for all $i, j \in \{1, \dots, n\}$ either $I_i = I_j$ or $I_i \cap I_j = \emptyset$.

Moreover, they have also the following nice property. Fix $w \in (K)^n$ and denote by $w_{U_j} = \{w_{\iota}, \text{ for any } \iota \in U_j\}$. Suppose we receive all the symbols in U_j . There is a simple linear parity test on the $r+1$ symbols of U_j such that if this parity check fails we know that at least one of the symbols in U_j is wrong. If we are guaranteed (or we assume) that at most one of the symbols in U_j is wrong and the parity check is OK, then all the symbols in U_j are correct. Moreover we can recover an erased symbol w_{ι} , with $\iota \in U_j$ using a polynomial interpolation through the points of the recovering set w_{U_j} .

3 GENERALIZED HAMMING WEIGHTS OF AG LRC CODES

Let K be a field and let \mathcal{X} be a smooth and geometrically connected curve of genus $g \geq 2$ defined over the field K . We also assume $\mathcal{X}(K) \neq \emptyset$. We recall the following definitions:

Definition 3 ([29], [30]). The **K-gonality** $\gamma_K(\mathcal{X})$ of \mathcal{X} over a field K is the smallest possible degree of a dominant rational map $\mathcal{X} \rightarrow \mathbb{P}_K^1$. For any field extension L of K , we define also the **L-gonality** $\gamma_L(\mathcal{X})$ of \mathcal{X} as the gonality of the base extension $\mathcal{X}_L = \mathcal{X} \times_K L$. It is an invariant of the function field $L(\mathcal{X})$ of \mathcal{X}_L .

Moreover, for each integer $i > 0$, the *i-th gonality* $\gamma_{i,L}(\mathcal{X})$ of \mathcal{X} is the minimal degree z such that there is $R \in \text{Pic}^z(\mathcal{X})(L)$ with $h^0(R) \geq i+1$. The sequence $\gamma_{i,\bar{K}}(\mathcal{X})$ is the usual *gonality sequence* [20]. Moreover, the integer $\gamma_{1,K}(\mathcal{X}) = \gamma_K(\mathcal{X})$ is the K-gonality of \mathcal{X} .

Let $K = \mathbb{F}_q$ a finite field with q elements. Let $C \subset K^n$ be a linear $[n, k]$ code over K . We recall that the *support* of C is defined as follows

$$\text{supp}(C) = \{i \mid c_i \neq 0 \text{ for some } c \in C\}.$$

So $\#\text{supp}(C)$ is the number of nonzero columns in a generator matrix for C . Moreover, for any $1 \leq v \leq k$, the *v-th generalized Hamming weight* of C [14, §7.10], [36, §1.1] is defined by

$$d_v(C) = \min\{\#\text{supp}(\mathcal{D}) \mid \mathcal{D} \text{ is a linear subcode of } C \text{ with } \dim(\mathcal{D}) = v\}.$$

In other words, for any integer $1 \leq v \leq k$, $d_v(C)$ is the v -th minimum support weights, i.e. the minimal integer t such that there are an $[n, v]$ subcode \mathcal{D} of C and a subset $S \subset \{1, \dots, n\}$ such that $\#(S) = t$ and each codeword of \mathcal{D} has zero coordinates outside S . The sequence $d_1(C), \dots, d_k(C)$ of generalized Hamming weights (also called *weight hierarchy* of C) is strictly increasing (see Theorem 7.10.1 of [14]). Note that $d_1(C)$ is the minimum distance of the code C .

Let us consider \mathcal{X} and \mathcal{Y} smooth projective absolutely irreducible curves over K and let $g : \mathcal{X} \rightarrow \mathcal{Y}$ be a rational separable map of curves of degree $r + 1$. Moreover we take $r, t, Q_\infty, f_1, \dots, f_m$ and $\mathcal{A} = g^{-1}(S)$ defined as Section 2.2. So we can construct an (n, k, r) algebraic geometric LRC code C as in Theorem 1. For this code we have the following:

Theorem 2. *Let C be an (n, k, r) algebraic geometric LRC code as in Theorem 1. For every integer $v \geq 2$ we have that*

$$d_v(C) \geq n - t\ell(r + 1) - (r - 1)h + \gamma_{v-1, K}(\mathcal{X}).$$

Proof. Take a v -dimensional linear subspace \mathcal{D} of C and call

$$E \subseteq \{P_{ij} \mid i = 0, \dots, r, j = 1, \dots, s\},$$

the set of common zeros of all elements of \mathcal{D} . Since $n - d_v(C) = \#(E)$, we have to prove that $t\ell(r + 1) + (r - 1)h - \#(E) \geq \gamma_{v-1, K}(\mathcal{X})$. Fix $u \in \mathcal{D} \setminus \{0\}$ and let F_u denote the zeros of u . Note that F_u is contained in the set $\{P_{ij} \mid i = 0, \dots, r, j = 1, \dots, s\}$ by the definition of the code C . We have $F_u \supseteq E$. By the definition of the integers t, ℓ and $h := \deg(x)$, we have $\#(F_u) \leq t\ell(r + 1) + (r - 1)h$. The divisors $F_u - E$, $u \in \mathcal{D} \setminus \{0\}$ form a family of linearly equivalent non-negative divisors, each of them defined over K . Since $\dim(\mathcal{D}) = v$, the definition of $\gamma_{v-1, \bar{K}}(\mathcal{X})$ gives $\#(F_u) - \#(E) \geq \gamma_{v-1, K}(\mathcal{X})$. This inequality for a single $u \in \mathcal{D} \setminus \{0\}$ proves the theorem. \square

See Remark 1 for an application of Theorem 2.

4 LRC CODES FROM NORM-TRACE CURVE

In this section we propose a new family of Algebraic Geometric LRC codes, that is, a LRC codes from the Norm-Trace curve. Moreover, we compute the \mathbb{F}_{q^u} -gonality of the Norm-Trace curve.

Let $K = \mathbb{F}_{q^u}$ be a finite field, where q is a power of a prime. We consider the *norm* $N_{\mathbb{F}_q}^{\mathbb{F}_{q^u}}$ and the *trace* $\text{Tr}_{\mathbb{F}_q}^{\mathbb{F}_{q^u}}$, two functions from \mathbb{F}_{q^u} to \mathbb{F}_q defined as

$$N_{\mathbb{F}_q}^{\mathbb{F}_{q^u}}(x) = x^{1+q+\dots+q^{u-1}} \text{ and } \text{Tr}_{\mathbb{F}_q}^{\mathbb{F}_{q^u}}(x) = x + x^q + \dots + x^{q^{u-1}}.$$

The *Norm-Trace curve* χ is the curve defined over K by the following affine equation

$$N_{\mathbb{F}_q}^{\mathbb{F}_{q^u}}(x) = \text{Tr}_{\mathbb{F}_q}^{\mathbb{F}_{q^u}}(y),$$

that is,

$$x^{(q^u-1)/(q-1)} = y^{q^{u-1}} + y^{q^{u-2}} + \dots + y \text{ where } x, y \in K \quad (1)$$

The Norm-Trace curve χ has exactly $n = q^{2u-1}$ K -rational affine points (see Appendix A of [5]), that we denote by $\mathcal{P}_\chi = \{P_1, \dots, P_n\}$. The genus of χ is $g = \frac{1}{2}(q^{u-1} - 1)(\frac{q^u-1}{q-1} - 1)$. Note that if we consider $u = 2$, we obtain the Hermitian curve.

Starting from the Norm-Trace curve, we have two different ways to construct Norm-Trace LRC codes.

PROJECTION ON x We have to construct a q^u -ary (n, k, r) LRC codes. We consider the natural projection $g(x, y) = x$. Then the degree of g is $q^{u-1} = r + 1$ and the degree of y is $h = 1 + q + \dots + q^{u-1}$.

To construct the codes we consider $S = \mathbb{F}_{q^u}$ and $D = tQ_\infty$ for some $t \geq 1$. Then, using a construction of Theorem 1 we find the parameters for these Norm-Trace LRC codes.

Proposition 1. *A family of Norm-Trace LRC codes has the following parameters:*

$$n = q^{2u-1}, \quad k = mr = (t + 1)(q^{u-1} - 1)$$

and

$$d \geq n - tq^{u-1} - (q^{u-1} - 1)(1 + q + \dots + q^{u-1}).$$

PROJECTION ON y We have to construct a q^u -ary (n, k, r) LRC codes. We consider the other natural projection $g'(x, y) = y$. Then $\deg(g') = 1 + q + \dots + q^{u-1} = r + 1$. In this case we take $S = \mathbb{F}_{q^u} \setminus M$, where

$$M = \{a \in \mathbb{F}_{q^u} \mid a^{q^{u-1}} + a^{q^{u-2}} + \dots + a = 0\},$$

so $r = q + \dots + q^{u-1}$ and $h = \deg(x) = q^{u-1}$. Then, using Theorem 1 we have the following

Proposition 2. *A family of Norm-Trace LRC codes has the following parameters:*

$$n = q^{2u-1} - q^{u-1}, \quad k = mr = (t + 1)(q + \dots + q^{u-1})$$

and

$$d \geq n - tq^{u-1} - (q + \dots + q^{u-1}) - q^{u-1}(q^{u-1} + \dots + q - 1).$$

For the Norm-Trace curve χ we are able to find the K -gonality of χ .

Lemma 1. *Let χ be a Norm-Trace curve defined over \mathbb{F}_{q^u} , where $u \geq 2$. We have $\gamma_{1, \mathbb{F}_{q^u}}(\chi) = q^{u-1}$.*

Proof. The linear projection onto the x axis has degree q^{u-1} and it is defined over \mathbb{F}_q and hence over \mathbb{F}_{q^u} . Thus $\gamma_{1, \mathbb{F}_{q^u}}(\chi) \leq q^{u-1}$. Denote by $z = \gamma_{1, \mathbb{F}_{q^u}}(\chi)$ and assume that $z \leq q^{u-1} - 1$. By the definition of K -gonality, there is a non-constant morphism $w: \chi \rightarrow \mathbb{P}^1$ with $\deg(w) = z$ and defined over \mathbb{F}_{q^u} . Since $w(\chi(\mathbb{F}_{q^u})) \subseteq \mathbb{P}^1(\mathbb{F}_{q^u})$, we get $\sharp(\chi(\mathbb{F}_{q^u})) \leq z(q^u + 1) \leq (q^{u-1} - 1)(q^u + 1)$, that is a contradiction. \square

Remark 1. By Lemma 1, we can apply Theorem 2 to the Norm-Trace curve. In fact, we can consider the gonality sequence over K of χ to get a lower bound on the second generalized Hamming weight of the two families of Norm-Trace LRC codes:

- Let $t \geq 1$ and let C be a $(q^{2u-1}, (t+1)(q^{u-1}-1), q^{u-1}-1)$ Norm-Trace LRC code. Then we have

$$d_2(C) \geq q^{2u-1} + q^{u-1} - tq^{u-1} - (q^{u-1}-1)(1+q+\dots+q^{u-1}).$$

- Let $t \geq 1$ and let C be a Norm-Trace LRC code with parameters $(q^{2u-1}-q^{u-1}, (t+1)(q+\dots+q^{u-1}), q+\dots+q^{u-1})$. Then we have

$$d_2(C) \geq q^{2u-1} - (t-1)q^{u-1} - (1+q^{u-1})(q+\dots+q^{u-1}).$$

5 HERMITIAN LRC CODES

In this section we improve the bound on the distance of Hermitian LRC codes proposed in [1] using some properties of *Hermitian codes* which are a special case of algebraic geometric codes.

5.1 Hermitian codes

Let us consider $K = \mathbb{F}_{q^2}$ a finite field with q^2 elements. The *Hermitian curve* \mathcal{H} is defined over K by the affine equation

$$x^{q+1} = y^q + y \text{ where } x, y \in K. \quad (2)$$

This curve has genus $g = \frac{q(q-1)}{2}$ and has $q^3 + 1$ points of degree one, namely a pole Q_∞ and $n = q^3$ rational affine points, denoted by $\mathcal{P}_{\mathcal{H}} = \{P_1, \dots, P_n\}$ [31].

Definition 4. Let $m \in \mathbb{N}$ such that $0 \leq m \leq q^3 + q^2 - q - 2$. Then the **Hermitian code** $C(m, q)$ is the code $C(D, mQ_\infty)$ where

$$D = \sum_{\alpha^{q+1} = \beta^q + \beta} P_{\alpha, \beta}$$

is the sum of all places of degree one (except Q_∞ , that is a point at infinity) of the Hermitian function field $K(\mathcal{H})$.

By Lemma 6.4.4. of [33] we have that

$$\mathcal{B}_{m, q} = \{x^i y^j \mid qi + (q+1)j \leq m, 0 \leq i \leq q^2 - 1, 0 \leq j \leq q - 1\},$$

forms a basis of $\mathcal{L}(mQ_\infty)$. For this reason, the Hermitian code $C(m, q)$ could be seen as $\text{Span}_{\mathbb{F}_{q^2}} \langle \text{ev}_{\mathcal{P}_{\mathcal{H}}}(\mathcal{B}_{m, q}) \rangle$. Moreover, the dual of $C(m, q)$ denoted by $C(m_\perp, q) = C^\perp(m, q)$ is again an Hermitian code and it is well known (Proposition 8.3.2 of [33]) that the degree m of the divisor has the following relation with respect to m_\perp :

$$m_\perp = n + 2g - 2 - m. \quad (3)$$

The Hermitian codes can be divided in four phases [11], any of them having specific explicit formulas linking their dimension and their distance [22]. In particular we are interested in the first and the last phase of Hermitian codes, which are:

I PHASE: $0 \leq m_{\perp} \leq q^2 - 2$. Then we have $m_{\perp} = aq + b$ where $0 \leq b \leq a \leq q - 1$ and $b \neq q - 1$. In this case, the distance is

$$\begin{cases} d = a + 1 & \text{if } a > b \\ d = a + 2 & \text{if } a = b. \end{cases} \quad (4)$$

IV PHASE: $n - 1 \leq m_{\perp} \leq n + 2g - 2$. In this case $m_{\perp} = n + 2g - 2 - aq - b$ where a, b are integers such that $0 \leq b \leq a \leq q - 2$ and the distance is

$$d = n - aq - b. \quad (5)$$

5.2 Bound on distance of Hermitian LRC codes

Let $K = \mathbb{F}_{q^2}$ be a finite field, where q is a power of a prime. Let $\mathcal{X} = \mathcal{H}$ be the Hermitian curve with affine equation as in (2). We recall that this curve has $q^3 \mathbb{F}_{q^2}$ -rational affine points plus one at infinity, that we denoted by Q_{∞} .

We consider two of the three constructions of Hermitian LRC codes proposed in [1] and we improve the bound on distance of Hermitian LRC codes using properties of Hermitian codes. In particular, if we find an Hermitian code $C(m, q) = C_{\text{Her}}$ such that $C_{\text{LRC}} \subset C_{\text{Her}}$, then we have $d_{\text{LRC}} \geq d_{\text{Her}}$.

PROJECTION ON \mathcal{X} By Proposition 4 of [1], we have a family of (n, k, r) Hermitian LRC codes with $r = q - 1$, length $n = q^3$, dimension $k = (t - 1)(q - 1)$ and distance $d \geq n - tq - (q - 2)(q + 1)$. Moreover, for these codes, $S = K$, $D = tQ_{\infty}$ for some $1 \leq t \leq q^2 - 1$ and the basis for the vector space V is

$$\mathcal{B} = \{x^j y^i \mid j = 0, \dots, t, i = 0, \dots, q - 2\}. \quad (6)$$

Using the Hermitian codes, we improve the bound on the distance for any integer t , such that $q^2 - q + 1 \leq t \leq q^2 - 1$.

To find an Hermitian code $C(m, q) = C_{\text{Her}}$ such that $C_{\text{LRC}} \subset C_{\text{Her}}$, we have to compute the set $\mathcal{B}_{m, q}$, that is, we have to find m . After that, to compute the distance of $C(m, q)$ we use (4) and (5).

We consider the first Hermitian phase: $0 \leq m_{\perp} \leq q^2 - 2$, that is, $q^2 - q + 1 \leq t \leq q^2 - 1$.

For this phase $m_{\perp} = aq + b$, where $0 \leq b \leq a \leq q - 1$ and the distance of the Hermitian code is either $d = a + 1$ if $a > b$ or $d = a + 2$ if $a = b$. By (6), m must be equal to $m = qt + (q + 1)(q - 2)$ and by (3) we have that $m_{\perp} = n + 2g - 2 - m = q(q^2 - t)$. So $b = 0$ and $a = q^2 - t$ and the distance of the Hermitian code is $d_{\text{Her}} = a + 1 = q^2 - t + 1$, since $a > b$. This implies that

$$d_{\text{LRC}} \geq q^2 - t + 1, \text{ for any } t \geq q^2 - q + 1. \quad (7)$$

Note that (7) improves the bound on the distance proposed in Proposition 4 of [1] since

$$q^2 - t + 1 > q^3 - tq - (q - 2)(q + 1) \iff t(q - 1) > q(q - 1)^2 + 1 \iff t > q^2 - q.$$

We just proved the following:

Proposition 3. *Let $q^2 - q + 1 \leq t \leq q^2 - 1$. It is possible to construct a family of (n, k, r) Hermitian LRC codes $\{C_t\}_{q^2 - q + 1 \leq t \leq q^2 - 1}$ with the following parameters:*

$$n = q^3, k = (t - 1)(q - 1), r = q - 1 \text{ and } d \geq q^2 - t + 1.$$

TWO RECOVERING SETS In [1] the authors propose an Hermitian code with two recovering sets of size $r_1 = q - 1$ and $r_2 = q$, denoted by LRC(2). They consider

$$L = \text{Span}\{x^i y^j, i = 0, \dots, q - 2, j = 0, \dots, q - 1\}$$

and a linear code C obtained by evaluating the functions in L at the points of $B = g^{-1}(\mathbb{F}_{q^2} \setminus M)$, where $g(x, y) = x$ and $M = \{a \in \mathbb{F}_q \mid a^q + a = 0\}$. So $|B| = q^3 - q$. By Proposition 4.3 of [1], the LRC(2) code has length $n = (q^2 - 1)q$, dimension $k = (q - 1)q$ and distance

$$d \geq (q + 1)(q^2 - 3q + 3) = q^3 - 2q^2 + 3. \quad (8)$$

As before, we improve the bound on the distance using Hermitian codes that contains the LRC(2) code. To do this we have to find m_\perp . By L , we have that $m = q(q - 1) + (q + 1)(q - 2)$ so we are in the fourth phase of Hermitian codes because $m_\perp = n + 2g - 2 - m = q^3 - q^2 + q$. In this case $d_{\text{Her}} = m_\perp - 2g + 2 = q^3 + 2q + 2$. Since $|B| = q^3 - q$, we have that

$$d_{\text{LRC}} \geq d_{\text{Her}} - q = q^3 + q + 2. \quad (9)$$

Note that this bound improves bound (8). We just proved the following proposition:

Proposition 4. *Let C be a linear code obtained by evaluating the functions in L at the points of B . Then C has the following parameters:*

$$n = (q^2 - 1)q, k = (q - 1)q, r_1 = q - 1, r_2 = q \text{ and } d \geq q^3 + q + 2.$$

ACKNOWLEDGEMENTS

The authors would like to thank the anonymous referees for their comments.

REFERENCES

- [1] A. Barg, I. Tamo, and S. Vlăduț. Locally recoverable codes on algebraic curves. *arXiv preprint arXiv:1501.04904*, 2015.
- [2] H. Chen, R. Cramer, S. Goldwasser, R. De Haan, and V. Vaikuntanathan. Secure computation from random error correcting codes. In *Advances in Cryptology-EUROCRYPT 2007*, pages 291–310. Springer, 2007.

- [3] M. Forbes and S. Yekhanin. On the locality of codeword symbols in non-linear codes. *Discrete Mathematics*, 324:78–84, 2014.
- [4] G. D. Forney. Dimension/length profiles and trellis complexity of linear block codes. *Information Theory, IEEE Transactions on*, 40(6):1741–1752, 1994.
- [5] O. Geil. On codes from Norm-Trace curves. *Finite Fields Appl.*, 9:351–371, 2003.
- [6] O. Geil, S. Martin, R. Matsumoto, D. Ruano, and Y. Luo. Relative generalized Hamming weights of one-point algebraic geometric codes. *Information Theory, IEEE Transactions on*, 60(10):5938–5949, 2014.
- [7] P. Gopalan, V. Guruswami, and P. Raghavendra. List decoding tensor products and interleaved codes. *SIAM Journal on Computing*, 40(5):1432–1462, 2011.
- [8] P. Gopalan, C. Huang, H. Simitci, and S. Yekhanin. On the locality of codeword symbols. *Information Theory, IEEE Transactions on*, 58(11):6925–6934, 2012.
- [9] V. Guruswami. List decoding from erasures: Bounds and code constructions. *Information Theory, IEEE Transactions on*, 49(11):2826–2833, 2003.
- [10] P. Heijnen and R. Pellikaan. Generalized Hamming weights of q -ary Reed-Muller codes. In *IEEE Trans. Inform. Theory*. Citeseer, 1998.
- [11] T. Høholdt, J. H. van Lint, and R. Pellikaan. Algebraic geometry of codes. In V. S. Pless and W. Huffman, editors, *Handbook of coding theory*, Vol. I, II, pages 871–961. North-Holland, 1998.
- [12] M. Homma and S. J. Kim. The second generalized hamming weight for two-point codes on a Hermitian curve. *Designs, Codes and Cryptography*, 50(1):1–40, 2009.
- [13] C. Huang, M. Chen, and J. Li. Pyramid codes: Flexible schemes to trade space for access efficiency in reliable data storage systems. *ACM Transactions on Storage (TOS)*, 9(1):3, 2013.
- [14] W. C. Huffman and V. Pless. *Fundamentals of error-correcting codes*. Cambridge university press, 2003.
- [15] H. Janwa and A. K. Lal. On generalized Hamming weights and the covering radius of linear codes. In *Applied algebra, algebraic algorithms and error-correcting codes*, pages 347–356. Springer, 2007.
- [16] T. Kasami, T. Takata, T. Fujiwara, and S. Lin. On complexity of trellis structure of linear block codes. *Information Theory, IEEE Transactions on*, 39(3):1057–1064, 1993.
- [17] T. Kasami, T. Takata, T. Fujiwara, and S. Lin. On the optimum bit orders with respect to the state complexity of trellis diagrams for binary linear codes. *Information Theory, IEEE Transactions on*, 39(1):242–245, 1993.

- [18] J. Kurihara and T. Uyematsu. Strongly-secure secret sharing based on linear codes can be characterized by generalized Hamming weight. In *Communication, Control, and Computing (Allerton), 2011 49th Annual Allerton Conference on*, pages 951–957. IEEE, 2011.
- [19] J. Kurihara, T. Uyematsu, and R. Matsumoto. Secret sharing schemes based on linear codes can be precisely characterized by the relative generalized hamming weight. *IEICE Transactions on Fundamentals of Electronics, Comications and Computer Sciences*, 95(11):2067–2075, 2012.
- [20] H. Lange and G. Martens. On the gonality sequence of an algebraic curve. *Manuscripta mathematica*, 137(3-4):457–473, 2012.
- [21] Y. Luo, C. Mitrpant, A. J. H. Vinck, and K. Chen. Some new characters on the wire-tap channel of type II. *Information Theory, IEEE Transactions on*, 51(3):1222–1229, 2005.
- [22] C. Marcolla. *On structure and decoding of Hermitian codes*. PhD thesis, University of Trento, 2013.
- [23] S. Martin and O. Geil. Relative generalized Hamming weights of q-ary Reed-Muller codes. *arXiv preprint arXiv:1407.6185*, 2014.
- [24] C. Munuera. On the generalized Hamming weights of geometric Goppa codes. *IEEE transactions on information theory*, 40(6):2092–2099, 1994.
- [25] C. Munuera and D. Ramirez. The second and third generalized Hamming weights of Hermitian codes. *Information Theory, IEEE Transactions on*, 45(2):709–712, 1999.
- [26] C.-K. Ngai, R. W. Yeung, and Z. Zhang. Network generalized hamming weight. *Information Theory, IEEE Transactions on*, 57(2):1136–1143, 2011.
- [27] W. Olaya-León and C. Granados-Pinzón. The second generalized hamming weight of certain Castle codes. *Designs, Codes and Cryptography*, 76(1):81–87, 2015.
- [28] D. S. Papailiopoulos and A. G. Dimakis. Locally repairable codes. *Information Theory, IEEE Transactions on*, 60(10):5843–5855, 2014.
- [29] R. Pellikaan. On the gonality of curves, abundant codes and decoding. In *Coding theory and algebraic geometry*, pages 132–144. Springer, 1992.
- [30] B. Poonen. Gonality of modular curves in characteristic p. *Mathematical Research Letters*, 14(4):691–701, 2007.
- [31] H. G. Ruck and H. Stichtenoth. A characterization of Hermitian function fields over finite fields. *Journal fur die Reine und Angewandte Mathematik*, 457:185–188, 1994.
- [32] N. Silberstein, A. S. Rawat, O. O. Koyluoglu, and S. Vishwanath. Optimal locally repairable codes via rank-metric codes. In *Information Theory Proceedings (ISIT), 2013 IEEE International Symposium on*, pages 1819–1823. IEEE, 2013.

- [33] H. Stichtenoth. *Algebraic function fields and codes*. Universitext. Springer-Verlag, Berlin, 1993.
- [34] I. Tamo and A. Barg. A family of optimal locally recoverable codes. *Information Theory, IEEE Transactions on*, 60(8):4661–4676, 2014.
- [35] I. Tamo, D. S. Papailiopoulos, and A. G. Dimakis. Optimal locally repairable codes and connections to matroid theory. In *Information Theory Proceedings (ISIT), 2013 IEEE International Symposium on*, pages 1814–1818. IEEE, 2013.
- [36] M. Tsfasman, S. Vlăduț, and D. Nogin. *Algebraic geometric codes: basic notions*, volume 139. American Mathematical Soc., 1990.
- [37] V. K. Wei. Generalized Hamming weights for linear codes. *Information Theory, IEEE Transactions on*, 37(5):1412–1418, 1991.
- [38] K. Yang, P. V. Kumar, and H. Stichtenoth. On the weight hierarchy of geometric Goppa codes. *Information Theory, IEEE Transactions on*, 40(3):913–920, 1994.