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(Article begins on next page)
Best Responding to What? A Behavioural Approach to One-Shot Play in 2 × 2 Games*

Andrea Gallice

8.1 Introduction

How should an inexperienced agent who is facing a simple strategic situation, such as for instance a one-shot 2 × 2 game, decide which strategy to play? A rational, or Bayesian, approach for dealing with such a decision under uncertainty (in a spirit similar to Savage, 1954) would be the following:

1 – the player forms a belief about what his opponent will play.
2 – the player chooses the strategy which best responds to this belief.

Interpreting it as a heuristics and acknowledging players' heterogeneity, this rule cannot be expected to describe the behaviour of every single agent. Still some recent papers confirm that in simple strategic interactions the majority of individuals behave in a manner consistent with their beliefs. Nyarko and Schotter (2002) study a 60-times repeated 2 × 2 game and find that around 75 per cent of the players do indeed best respond to their stated beliefs. For the case of 3 × 3 games, Rey-Biel (2004) considers 10 one-shot games and finds a similar rate of compliance while 55 per cent is the percentage found by Costa-Gomes and Weizsäcker (2005) using data about 14 (more complex) one-shot games. It is therefore a conservative guess to expect that, in one-shot 2 × 2 games, at least half of the individuals play consistently with their beliefs. We want to capture the behaviour of this majority of players.

In order to do so we focus initially on the process of beliefs formation. The goal is to find a function that may approximate players' beliefs without the need of having to elicit them explicitly. Using an axiomatic approach, we first list a limited number of desirable properties that, in accordance with experimental studies and behavioural regularities, a belief function must fulfill. Then, we check some existing concepts commonly used in game theory and decision theory to see which of them, if any, fulfills all the requirements.

We first show that, contrary to a very authoritative strand of research (starting with Aumann, 1987, who simplified Harary, 1973), the concept of mixed strategy Nash equilibrium is not a suitable candidate for approximating players' beliefs. Indeed mixed equilibria are still problematic both in theory and in practice (see the literature review). In particular some features of mixed equilibria clearly contradict our axiomatic description.

At the opposite we find minimax regret (a criterion for dealing with choices under uncertainty originally proposed by Savage, 1951) to be the unique candidate to obey all the axioms. Therefore we propose minimax regret as a proxy for players' beliefs and we claim that the majority of players play 'as if' they were best responding to these approximated beliefs.

This conjecture is tested in the second part of the chapter. The predictions stemming from the procedure (best respond to beliefs equal to the minimax regret distribution of the opponent) are compared with experimental evidences about different versions of 2 × 2 one-shot matching pennies games. To forecast players' choices in this class of games can be particularly problematic because the Nash indication is often misleading (see, for instance, Ochs, 1995 and Goeree and Holt, 2001). Our procedure proves to be an effective way to identify the strategies which are more likely to be played. In fact, it correctly predicts the actual choices of around 80 per cent of the players.

We also apply the procedure to other classes of 2 × 2 games and we analyse its relationship with the Nash prediction. An interesting result is that the procedure selects a single outcome even in games that have multiple Nash equilibria such that it contributes to the debate on equilibrium selection (see Straub, 1995 and Haruvy and Stahl, 2004).

A number of studies that focus on how people play one-shot simultaneous games and investigate the issue of beliefs formation are closely related to this chapter. For instance, Stahl and Wilson (1995), Costa-Gomes et al. (2000) and Camerer et al. (2004) hypothesize and test the existence and relative importance of various archetypes of players that differ in the prior they have about the degree of sophistication of their opponents. An important difference with respect to these papers is that we do not consider heterogeneity. We recognize heterogeneity to be a very important feature of human behaviour but here we still favour simplicity. Therefore we only model the behaviour of the majority of players. Our contribution is to provide a fast, easy to implement and effective procedure that gives rule of thumb predictions about the outcome one should expect to arise in simple games.

The chapter is structured as follows: section 8.2 briefly reviews the theoretical interpretations and the empirical relevance of mixed strategy Nash equilibria. Section 8.3 presents the basic axioms that must be satisfied by a

* The author would like to thank the editors of the book and an anonymous referee for useful comments as well as the Collegio Carlo Alberto for kind hospitality during the period in which the paper has been written.
beliefs function. Section 8.4 checks the compliance to these axioms of various candidate functions and it shows that minimax regret is the unique one to obey them all. Section 8.5 formally defines the procedure to be used to predict the outcome of simple games. These predictions are then tested in section 8.6. Section 8.7 concludes.

8.2 Literature review on mixed equilibria

The idea that agents may randomize over a set of actions (that is, they may use mixed strategies) dates back to Borel (1921) and was then enriched and developed by Von Neumann and Morgenstern (1944). The concept of Nash equilibrium appeared a few years after (Nash, 1951) and rapidly became the most important solution concept of game theory. In his paper Nash also presented the concept of mixed strategy Nash equilibrium, a notion that led to the famous theorem about the existence of a Nash equilibrium in any finite game.

![Diagram of the equilibrium problem](image)

**Figure 8.1** The equilibrium problem

In a mixed strategy Nash equilibrium (shortened in what follows to MSNE) each player randomizes over (some of) his pure strategies according to a probability distribution that makes the other players indifferent about what to play. Therefore no player has any strict incentive to deviate and the mixed strategy profile identifies a Nash equilibrium of the game.

The problem to look for this specific probability distribution shares quite many features with the following problem of elementary physics.

**Exercise:** You have a rigid rod of negligible weight which is 1 metre long. On one side of it there is a weight of 1 kg, on the other side a weight of 4 kg. Where does the rod have to be pivoted in order for it to be in equilibrium?

**Solution:** The rod is in equilibrium if $4l_1 = l_2$ where $l_1 + l_2 = 1$ and $l_1$ is the distance from the pivot to the weight of 4 kg and $l_2$ is the distance from the pivot to the weight of 1 kg. Therefore $l_1 = 0.2 \text{ m}$ and $l_2 = 0.8 \text{ m}$. The pivot has to be put closer to the heavier side. Indeed $\frac{3}{4}$ of the rod have to be given to the lighter side in order to impose the equilibrium.

A similar situation arises when we compute MSNE: the equilibrium condition among pure strategies is imposed and then, as a consequence, the appropriate probability distribution is retrieved. And as it happens in the rod example, mixed equilibria often allocate larger probabilities to ‘lighter’ strategies, i.e. to strategies that are associated with lower payoffs. This feature undermines the predictive power of MSNE. Another unappealing feature of mixed equilibria is the so-called ‘no own payoff effect’ which will be shortly discussed.

To sum up, MSNE often appear as quite an artificial construction and indeed they are still problematic both in theory and in practice. To quote Rubinstein (1991, p. 912): *The concept of mixed strategy has often come under heavy fire.* On the theoretical side different interpretations have been given about mixed equilibria and a general consensus is still missing. On the empirical side the relevance of MSNE in capturing agents’ behaviour has also been heavily questioned.

8.2.1 Theoretical interpretations of mixed equilibria

Taken at face value a mixed strategy Nash equilibrium prescribes a player to select the strategy to play according to a specific probability distribution. In other words, players should deliberately randomize. This view may make sense when players are interacting in repeated games and they do not want their pattern of play to be predictable. But in one-shot interactions such a randomizing behaviour looks less realistic.

A second interpretation considers the MSNE probabilities as indicating the steady state frequencies of pure strategies when games are played in large populations (see, for instance, Rosenthal, 1979). Each player chooses a pure strategy but in the entire population the MSNE distribution should emerge. Again, this interpretation does not look appropriate in the context we are interested in, namely a one-shot interaction between two players.

Very much related with the ideas presented in this chapter are the interpretations that link MSNE with the beliefs of the players. Harsanyi (1973) presented the so-called purification interpretation of MSNE. The idea is that players play pure strategies. The twist is that each player’s choice is based on some private information. This means that player $i$ knows what to play but the information sets of $i$’s opponents are not enough precise to allow them to be sure about $i$’s choice. MSNE captures this uncertainty (or ignorance).

Aumann (1987) further simplified this idea. The claim (see also Aumann and Brandenburger, 1995; Rényi and Robson, 2003) is that, even in the absence of this small amount of private information, players are still unsure about the opponents’ moves. Therefore the probability distribution that the MSNE assigns to player $i$ can be directly interpreted as the beliefs all the other players hold about $i$’s choice. We will come back to this point in section 8.4.1.

8.2.2 Empirical relevance of mixed equilibria

Because of the existence of all of these different theoretical interpretations, it is no surprise that many experimental studies have been designed with
the aim of testing the empirical relevance of MSNE. In particular, researchers focused on the study of matching pennies (MP) games, i.e., games that have a unique Nash equilibrium which is in mixed strategies. The common design of these experiments consisted in letting subjects repeatedly play the same version of a MP game. Results show (for a detailed review, see also Camerer, 2003) that long-run frequencies of pure strategies are not far from the MSNE predictions. Still players' behaviour at an individual level is often significantly different from that indicated by the equilibrium.

A different, and less investigated, question is to study how agents behave in front of a single interaction. In a one-shot MP game players cannot learn over time and the incentive to maintain an unpredictable pattern of play simply does not exist. In these cases the MSNE prediction is useless. In fact, the equilibrium indicates a mixed distribution while the result of the game will be a specific outcome given by the intersection of two pure strategies. Our chapter is focused on one-shot games for three main reasons: first, we claim our theory to be able to capture the behaviour of inexperienced players; secondly, as just mentioned, one-shot individuals' play has been less investigated; third, we think that enough real life situations are more likely to be similar to one-off events rather than to repeated interactions.

8.3 An axiomatic approach to model players' beliefs in 2 × 2 games

In this chapter we present some axioms that, in our view, capture some very basic properties of what the beliefs of inexperienced and boundedly rational players should look like. These axioms are quite general and therefore, perhaps unfortunately, they do not characterize a single belief function. However, they are sufficient to discriminate clearly among various functions that are commonly used in game theory and in decision theory. In fact, in order to find a function that may approximate players' beliefs, we do not want to introduce new ad hoc formulas. At the opposite we rely on concepts that already exist and that are widely used, though possibly for different purposes. Given that the conjectured belief function will be used for predictive purposes, we perceive this to be a useful result.

Consider the following game, where player 1 ∈ {A, B} can choose between strategies H and T,

<table>
<thead>
<tr>
<th></th>
<th>Hb</th>
<th>Tb</th>
</tr>
</thead>
<tbody>
<tr>
<td>b</td>
<td>βHA</td>
<td>βTA</td>
</tr>
<tr>
<td>c</td>
<td>βHA</td>
<td>βTA</td>
</tr>
<tr>
<td>d</td>
<td>βHA</td>
<td>βTA</td>
</tr>
</tbody>
</table>

We model the beliefs player B holds about what player A will play. This is the reason why the payoff matrix is incomplete and only the payoffs of player A appear. Whilst keeping in mind the example of an inexperienced boundedly rational player, B's beliefs are considered to be just a function of player A's payoffs. This may seem to be a drastic requirement, but it can be seen as a simplification of the fact that B's beliefs are mainly a function of A's payoffs and then, as a secondary effect that we do not consider, B's beliefs may also be influenced by higher order speculations that involve B's payoffs.

The focus on B's beliefs implies no loss of generality given that a similar analysis can be done for what concerns A's beliefs. We indicate with

- βHA the beliefs player B hold about player A playing strategy H;
- βTA the beliefs player B hold about player A playing strategy T.

Finally we put no particular restriction on the values of the payoffs a, b, c, and d. According to us, player B's belief function must obey the following five axioms:

[A1] Functional form

βHA = βHA(a, b, c, d), βTA = βTA(a, b, c, d)

[A2] Consistency with probability distribution

βHA ≥ 0, βTA ≥ 0 and βHA + βTA = 1 for any a, b, c, d.

The first axiom simply formalizes what we have already mentioned, i.e. the fact that B's beliefs have to be a function of the payoffs of player A. This implies that player B realizes A does respond to changes in his own payoffs. This own payoffs effect is a very robust feature of games played in experiments (for clear evidences of this effect see, among others, Ochs, 1995 and Goeree and Holt, 2001). The second axiom requires the function to identify a proper probability distribution — that is, beliefs that are non-negative and that sum up to one. Notice that, because of the relation βTA + βHA = 1, a single belief is enough to define the entire distribution. Therefore, from now on, we focus on βHA.

[A3] Continuity

βHA is continuous in all its arguments.

[A4] Monotonicity

βHA is weakly increasing in a and b, βHA is weakly decreasing in c and d.

[A5] Consistency with rows or columns switch

Continuity (A3) is required since there are no evident reasons for B's beliefs to jump in a discrete way given small changes in the arguments of the function. The monotonicity axiom (A4) defines the sign of the already mentioned own payoff effect and it is in line with a large amount of experimental evidences (among others Ochs, 1995; Goeree and Holt, 2001; Goeree et al.,
2003). The axiom states that players believe their opponent to be more likely to play strategies that look better. In other words, if the payoffs associated with strategy $H_A$ increase then the probability that player $B$ assigns to the event of $A$ playing that strategy cannot decrease.

The last axiom indicates that the beliefs of player $B$ have to react consistently to the payoffs structure of the game and they do not depend on the labelling of the strategies. To better understand the meaning of this axiom consider the two games that are reported below. With respect to the original game (Game 1 above), Game 1' is such that the payoffs of the two rows have been switched while Game 1'' is such that the payoffs of the two columns have been switched. Axiom 5 requires that, whenever $a < c$ and $b < d$, the probability that $B$'s belief function allocates to the event of $A$ playing strategy $H_A$ in Game 1 must be equal to the probability that the same function allocates to the event of $A$ playing $T_A$ in Game 1'. In other words $\beta_{H_A} = \beta'_{H_A}$. At the opposite the belief function must allocate the same probabilities to the events of $A$ playing $H_A$ in Game 1 and in Game 1'' such that $\beta_{H_A} = \beta''_{H_A}$. In fact a column switch does not change the relative attractiveness of strategies $H_A$ and $T_A$ from player $A$'s point of view.

8.4 The candidate functions

In this section we present various candidate functions and check their compliance with the axioms. We first analyse a number of proposals connected with the concept of mixed strategy Nash equilibrium. We then turn our attention to the minimax regret - a criterion normally used in decision theory. In line with the axiomatic analysis of the previous section, we check how the following functions perform in approximating the beliefs player $B$ holds about what player $A$ will play in Game 1.

### 8.4.1 Mixed strategy Nash equilibrium of player A

As mentioned in the literature review, the interpretation that mixed strategy equilibria could capture players' beliefs dates back to Harsanyi (1973) and it was later simplified by Aumann (1987). According to Aumann's view (see also Aumann and Brandenburger, 1995 and Reny and Robson, 2004) the probability distribution that the MSNE attaches to player $i$ represents the common belief all the other players have about what $i$ will play.

In the context of our 2x2 game this boils down to the hypothesis that $A$'s equilibrium distribution captures $B$'s beliefs. This interpretation suffers from one fundamental drawback. By construction the MSNE of player $A$ is exclusively a function of the payoffs of player $B$ and it is totally unrelated with $A$'s payoff. In the literature this unappealing feature of mixed equilibria is called 'no own payoff effect'. To make things clear consider Game 2 where, with respect to Game 1, also the payoffs of player $B$ appear.

$$
\begin{array}{ccc}
\beta_{H_A} & H_A & x, b, y \\
\beta_{T_A} & T_A & c, w, d, z \\
\end{array}
$$

The distribution that the MSNE assigns to player $A$ is defined by the probability $\hat{p}$ that makes player $B$ indifferent between playing $H_B$ or $T_B$. Therefore $\hat{p}$ solves $px + (1 - \hat{p})w = py + (1 - \hat{p})z$ such that $\hat{p} = \frac{x - z + y - w}{x + y + z - w}$ and the equilibrium component of player $A$ is given by $(\hat{p}H_A + (1 - \hat{p})T_A)$. Interpreting this probability distribution as $B$'s beliefs we would have:

$$
\beta_{H_A} = \frac{z - w}{x - y + z - w}
$$

Notice the counter-intuitive implications of this proposal: $B$'s beliefs remain constant no matter how $A$'s payoffs may change. For instance $\beta_{H_A} = \hat{p}$ in the game above as well as in a similar game where the payoffs $a$ (or $b$, $c$, $d$) is substituted with, for example, $5a$. This proposal therefore fails both Axiom 1 and Axiom 4. Beliefs of this kind would not capture any own payoff effect. This is a serious limitation given that the existence and the importance of such an effect is testified by many experimental studies.

Notice furthermore that the predictive power of this proposal would be null. In fact any strategy (or combinations of strategies) in support of the MSNE of player $B$ is a best response to these conjectured beliefs.

### 8.4.2 Mixed strategy Nash equilibrium of player B

The analysis of the previous section may suggest the use of the MSNE of player $B$ as a way to approximate $B$'s beliefs. Indeed the probabilities implied by the MSNE of $B$ are as a function of $A$'s payoffs. More precisely the MSNE of $B$ is given by $(\hat{q}H_B + (1 - \hat{q})T_B)$ where $\hat{q}$ solves $qa + (1 - \hat{q})b = qy + (1 - \hat{q})z$ such that $\hat{q} = \frac{a - z + b - d}{a + c + d - b}$ and therefore, under this proposal, we would have:

$$
\beta_{H_B} = \frac{d - b}{a - c + d - b}
$$

The problem with this formulation is that beliefs of this kind do not obey the monotonicity requirement. In fact $\beta_{H_B}$ is increasing in $a$ and $b$ and decreasing in $c$ and $d$ while Axiom 4 requires the opposite behavior. Assume that payoffs $a$ and $b$ increase such as to make strategy $H_A$ more attractive from
A's point of view. The probability weight that the MSNE attaches to the corresponding strategy \( H_B \) will get smaller (the intuition is the same provided by the rod example presented in section 8.2). Therefore the mixed equilibrium of player B does capture the own payoff effect but with the wrong sign.

Trying to improve on the limits of the last proposal one may be tempted to approximate B's beliefs by switching the probabilities implied by the MSNE of B. In other words by setting \( \beta_{HA} = 1 - \tilde{q} \) such that:

\[
\beta_{HA} = \frac{a - c}{a - c + d - b}
\]

This proposal is a function of A's payoffs and it now satisfies the monotonicity axiom. Still it fails Axiom 5. In fact, following this conjecture, player B should keep the same beliefs (\( \beta_{TB} = \tilde{p}_{TB} \)) also in game 1' (see below) where again the MSNE of B is given by \( \tilde{q} = \frac{d - b}{a - c + d - b} \) such that \( 1 - \tilde{q} = \beta_{TB} \). This would violate A5 whenever \( a \neq c \) and \( b \neq d \).

8.4.3 A new proposal: the minimax regret

We now present a new and unusual candidate for approximating players' beliefs. This proposal is based on an instrumental use of the minimax regret criterion. More precisely we claim that the beliefs of player A about what player i will play can be approximated by the minimax regret of player i.

Minimax regret, originally proposed by Savage (1951), is a concept which found its main applications as a selection criterion in decision theory (starting with Minford, 1954). The minimax regret criterion prescribes a player who has to make a decision under uncertainty to choose the action that minimizes the maximum regret he may suffer. The regret of player i is defined as the difference between the best payoff i could have got if he knew what his opponent (another player or Nature) had played and the payoff player i actually got.

The first step to properly compute the minimax regret consists in building the regret matrix which captures these differences. In the context of Game 1 (reproduced below) one needs to know which is the largest payoff between a and c and between b and d. Let us assume that \( a > c \) and \( d > b \). In this case the regret matrix is given by \( R_1 \):

\[
\begin{array}{ccc|cc}
\beta_{HA} & \beta_{TB} & H_B & T_B & H_A & 0
\end{array}
\]

A strategy \( H_A \) attains minimax regret if \( a - c > d - b \) while strategy \( T_A \) attains minimax regret with a reversed relation. Taking this specification as a belief function would clearly be unsatisfactory given that such a proposal would fail the continuity axiom. The use of mixed strategies solves this problem.

The mixed minimax regret is defined by the probability distribution (identified by \( \tilde{p}, \tilde{q} \) where the index r indicates regret) that it equals the expected regret of the two strategies. This optimal \( \tilde{p} \) solves \( p_r(a - d) = (1 - p_r)(a - c) \) so that \( p_r = \frac{a - c}{a - c + d - c} \). According to the conjecture of this paper \( \beta_{HA} = \tilde{p} \) should hold and thus:

\[
\beta_{HA} = \frac{a - c}{a - c + d - b}
\]

Once again, referring to Game 1 above, this means that player B approximately believes player A will play strategy \( H_A \) with probability \( \frac{a - c}{a - c + d - b} \) and strategy \( T_A \) with probability \( \frac{d - b}{a - c + d - b} \).

It is easy to show that this candidate function obeys all the axioms. In fact, \( \beta_{HA} = \beta_{TB} \) (a,b,c,d) (A1), \( \beta_{HA} \geq 0, \tilde{p}_{TB} \geq 0 \) and \( \beta_{HA} + \beta_{TB} = 1 \) (A2), \( \beta_{HA} \) is continuous in its arguments (A3) and it also fulfills the monotonicity requirement (A4). Finally it also obeys Axiom 5. In fact, referring to games 1' and 1" presented before, we have that \( \tilde{p}_r = \frac{a - c}{a - c + d - c} \) in Game 1' such that \( \beta_{HA} = \tilde{p}_r \) and \( \tilde{p}_r = \frac{a - c}{a - c + d - c} \) in Game 1" such that \( \beta_{HA} = \tilde{p}_r \).

8.4.4 A summary and an example

Table 8.1 summarizes how the various proposals that we just discussed perform in matching the five axioms. Axioms are identified as: functional form (A1), consistency with probability distribution (A2), continuity (A3), monotonicity (A4) and consistency with rows or columns switch (A5).

Among the various candidates, the (mixed) minimax regret proposal is the unique one to satisfy all the axioms and thus the unique one to qualify for approximating players' beliefs. Before properly defining a procedure which is based on these conjectured minimax regret beliefs, we present a useful example. This example is meant to show how simple is the process of approximating beliefs using minimax regret, under lining once more the inadequacy of proposals connected with the MSNE concept.

Consider Game 3 where \( k \in (-\infty, \infty) \) such that the game encompasses the cases of a two person game (for \( k \in (-1, \infty) \)) and of a game with a dominant strategy (for \( k \in (-\infty, 1) \).

\[
\begin{array}{ccc|cc}
H_B & T_B & H_A & 0, 2 & 2, 0
\end{array}
\]

\[
\begin{array}{ccc|cc}
H_B & T_B & k, 1 & 0, 2
\end{array}
\]

\[
\begin{array}{ccc|cc}
H_B & T_B & k + 1, 0 & 0, 2
\end{array}
\]
For $k \in (-\infty, -1]$ the regret matrix is given by $R_1$. The minimax regret of player $A$ is given by $(\hat{p}_A H_A + (1 - \hat{p}_A) T_A)$ where $\hat{p}_A = \frac{1}{2 + k}$. Following our proposal $\hat{p}_A = \frac{1}{2 + k}$ captures $A$’s beliefs about $A$ playing strategy $H_A$. This function appears as the bold concave curve in Figure 8.2 which focuses on the beliefs of player $B$ about what $A$ will play.

For $k \in (-\infty, -1]$ the regret matrix is given by $R_2$ such that $\hat{p}_B = \beta_h = 0$, i.e. player $B$ attaches a null probability to the event of $A$ playing $H_A$. In Figure 8.2 this appears as the bold line that lies on the horizontal axis for $k \leq -1$.

The other two functions (dashed lines) depict, respectively, the probability that the mixed strategy Nash equilibrium assigns to player $A$ playing strategy $H_A(\frac{1}{2})$ and to player $B$ playing strategy $H_B(\frac{1}{2})$.

### Table 8.1 Compliance to the axioms of the candidate functions

<table>
<thead>
<tr>
<th>Candidates/axioms</th>
<th>A1</th>
<th>A2</th>
<th>A3</th>
<th>A4</th>
<th>A5</th>
</tr>
</thead>
<tbody>
<tr>
<td>MSNE of pl. $A$</td>
<td>Y</td>
<td>N</td>
<td>Y</td>
<td>N</td>
<td>Y</td>
</tr>
<tr>
<td>MSNE of pl. $B$</td>
<td>Y</td>
<td>Y</td>
<td>Y</td>
<td>N</td>
<td>N</td>
</tr>
<tr>
<td>$1-MSNE$ of pl. $B$</td>
<td>Y</td>
<td>Y</td>
<td>Y</td>
<td>N</td>
<td>N</td>
</tr>
<tr>
<td>Pure minimax regret of pl. $A$</td>
<td>Y</td>
<td>Y</td>
<td>N</td>
<td>Y</td>
<td>Y</td>
</tr>
<tr>
<td>Mixed minimax regret of pl. $A$</td>
<td>Y</td>
<td>Y</td>
<td>Y</td>
<td>Y</td>
<td>Y</td>
</tr>
</tbody>
</table>

**Figure 8.2** Beliefs approximation through minimax regret in Game 3

The procedure

1. Compute the minimax regret distribution for the two players and retrieve $\hat{p}_A$ and $\hat{p}_B$.
2. Assign the following beliefs to the two players:
   - $\hat{p}_A = (\hat{p}_B, (1 - \hat{p}_B))$
   - $\hat{p}_B = (\hat{p}_A, (1 - \hat{p}_A))$
3. Let the two players choose the strategy to play according to $BR_i(\beta_i)$:
   - $BR_A(\beta_A) = \begin{cases} 
   H_A & \text{iff } u_i(H_A(\beta_A)) > u_i(T_A(\beta_A)) \\
   T_A & \text{iff } u_i(H_A(\beta_A)) < u_i(T_A(\beta_A)) \\
   0.5(H_A + 0.5T_A) & \text{iff } u_i(H_A(\beta_A)) = u_i(T_A(\beta_A)) 
   \end{cases}$

The strategies selected by $BR_i(\beta_i)$ are the ones which have the largest probability to be played in a one-shot game or, equivalently, the ones which we would expect to be chosen with the highest frequency if the game is played in a large enough population. In particular, whenever $u_i(H_i(\beta_i)) \neq u_i(T_i(\beta_i))$ for any $i$ then every player has a single best response and the intersection of the two selected strategies indicates a single outcome of the game as the most likely one to arise.

The procedure thus provides a forecast in three simple steps: it is enough to compute the minimax regret, use its probability distributions to approximate intersect just once. The intersection happens for the unique $k$ (in this case $k = 1$, symmetric game) for which all the three functions reach a value of $\frac{1}{2}$. 

**8.5 The procedure**

We use the conjectured minimax regret beliefs as a starting point for a procedure that selects the strategies more likely to be played by inexperienced agents interacting in one-shot $2 \times 2$ games. In order to formalize this procedure, we consider a generic $2 \times 2$ game where $s_i = (H_i, T_i)$ and $u_i(s_i, s_j)$ are the strategy space and the payoffs of player $i$ with $i, j \in \{A, B\}$. The minimax regret distributions (one for each player) are given by:

$$((\hat{p}_A H_A + (1 - \hat{p}_A) T_A), (\hat{p}_B H_B + (1 - \hat{p}_B) T_B))$$

where $\hat{p}_i$ defines the probability with which player $i$ has to play strategy $H_i$ in order to minimize his regret.

With a slightly different notation with respect to the previous sections where only $B$’s beliefs were considered, define now as $\beta_i = (\theta, 1 - \theta)$ the beliefs of player $i$ such that $i$ believes that player $j$ will play strategies $H_j$ with probability $\theta$ and strategy $T_j$ with probability $1 - \theta$.

$BR_i(\beta_i)$ is the best reply function of player $i$. It uses $i$’s beliefs as an input and provides as an output the strategy that $i$ must choose in order to maximize his expected payoff.

The procedure

1. Compute the minimax regret distribution for the two players and retrieve $\hat{p}_A$ and $\hat{p}_B$.
2. Assign the following beliefs to the two players:
   - $\hat{p}_A = (\hat{p}_B, (1 - \hat{p}_B))$
   - $\hat{p}_B = (\hat{p}_A, (1 - \hat{p}_A))$
3. Let the two players choose the strategy to play according to $BR_i(\beta_i)$:
players' beliefs and choose for each player the strategy that best responds to these beliefs.\textsuperscript{13} We do not claim this procedure to be consciously used by players. What we claim is that, on average, the procedure is operationally valid, i.e. the majority of individuals play the game 'as if' they were applying it.

### 8.6 Experimental evidences about matching pennies games

We apply the proposed procedure to MP games for which experimental results are available from other studies.\textsuperscript{18} Given that the procedure aims to capture the behaviour of inexperienced players, we only consider experiments where subjects played a single game just once. The data are reported in Table 8.2.

The first three games (GH1, GH2 and GH3) and the correspondent experimental results are taken from Goeree and Holt (2001). Each game was played once by a different pool of 50 subjects. In the original paper the authors use these games to evaluate the predictive power of the mixed strategy Nash equilibrium. The last three games appear in Goeree and Holt (2004) who took them from Guyer and Rapoport (1972). In the original experiment 214 subjects played in a random order 244 games belonging to different typologies.

We now explain the meaning of the last four columns of Table 8.2. In the fourth to last column we report BR(\(p_i\)), the strategy selected by the procedure. The third to last column presents the experimental results in the form \(a/b S_i\), where \(a\) is the number of players that chose strategy \(S_i\) in \([H_i, T_i]\) and \(b = 0.5N\) is the total number of row or column players.

The second to last column shows the hit rate which measures the performance of the prediction in forecasting actual behaviour. The hit rate is a simple summary statistics which counts the number of hits: it ranges between 0 per cent (all misses) and 100 per cent (all hits).\textsuperscript{14} Therefore, when the procedure indicates a single strategy, the hit rate simply captures the percentage of players who actually played it. In games in which the procedure indicates that subjects should uniformly randomize and \(b\) is odd, the hit rate reaches 100 per cent if the players split as equally as possible. In game GH1 for instance the hit rate would have been 100 per cent both if 12 or 13 out of the 25 row or column players chose \(H_1\).

Finally in the last column of Table 8.2 we test for the validity of our conjecture—that is, we test the hypothesis of the procedure being able to \textit{ex-ante} predict the strategies that are overplayed. Therefore we only consider the cases in which the procedure selects a pure strategy. Using a one-sided test, we test if the proportion of players that plays \(BR(p_i)\) is significantly greater than 50 per cent. To do so we use Fisher's exact probability test which calculates the difference between the observed data and an alternative data distribution. When our procedure selects a pure strategy we expect the null hypothesis (observed data not being significantly different from the uniform distribution) to be rejected. In other words we expect the p-value that appears in the last column to be below 5 per cent. An underlined p-value indicates that this is the case.

To have a better feeling of how the procedure actually works consider a couple of examples. Game GH1 is a symmetric matching pennies game. The minimax regret is obviously \(\frac{1}{2}H_1 + \frac{1}{2}T_1\) for any \(i \in \{A, B\}\) and thus the procedure assigns uniform beliefs to both players. Both strategies lead to the same expected payoff and the procedure predicts all outcomes to be equally likely. Even if the population is quite small, actual frequencies confirm that the distributions of players' choices are as uniform as possible.

In game GH2 the minimax regret distributions are given by \(\{\frac{1}{4}H_1 + \frac{1}{4}T_1\} \frac{1}{2}H_1 + \frac{1}{2}T_1\). The procedure then selects strategies \(H_1\) and \(T_1\) with equal probability. This is the best response to the conjectured minimax regret beliefs.\textsuperscript{15} Strategy \(H_1\) was actually chosen by 24 out of the 25 row players (hit rate of 96 per cent) and strategy \(T_1\) was chosen by 21 of the 25 column players (84 per cent).

The hit rate remains above 75 per cent and the p-values are in line with our conjecture also in games GH3, GR4 and GR5 while results are less good in the case of Game GR6 where the procedure failed to predict that B players overplayed strategy \(H_4\).\textsuperscript{15} Considering only the games where the procedure indicates a single outcome (GH2, GH3, GR4, GR5), the procedure correctly predicts the choices of 81 per cent of the players.

It can be shown easily that predictions based on letting players best respond to beliefs approximated by the MSNE clearly underperform our procedure. We already mentioned the fact that using the MSNE of player \(i\) to approximate \(j\)'s beliefs (Aumann, 1987) does not provide any specific indication given that
all the strategies of \( j \) are a best response to these beliefs. And to use the MSNE of player \( i \) to approximate \( i \)'s beliefs often leads to misleading results (for instance, in Game GH2 the outcome selected would be \((H_A, H_B)\) instead of \((H_A, T_B)\)).

Moreover, also interpreting MSNE as a prediction in itself (such that, for instance, players \( A \) should uniformly randomize in all the three GH games given that his mixed equilibrium is always \( \frac{1}{3} H_A + \frac{2}{3} T_A \)) is not effective. In fact it works fine only in the case of symmetric games (like GH1) while it is clearly inadequate in the other games. Indeed Goeree and Holt (2001) present the results of Game GH1 as supportive of the MSNE prediction, while they show the results of games GH2 and GH3 as evidences of its failure. Therefore they write (p. 1419) that ‘The Nash analysis seems to work only by coincidence, when the payoff structure is symmetric and deviation risks are balanced.’

Analysing the same results through the lens of our conjecture, it seems indeed that the fact that the Nash analysis works in game GH1 is the result of a coincidence. But there is an explanation for this coincidence. In symmetric MP games the probability distributions implied by the MSNE and by the minimax regret always coincide.\(^{17}\) Still, as soon as the structure of the game becomes asymmetric, individuals’ behaviour is far better captured by our behavioural model rather than by the Nash prediction.

8.6.1. The procedure in other games

Until now we have only considered the more problematic case of MP games, but the procedure can be applied to any \( 2 \times 2 \) one-shot game. The steps to select the strategies that inexperienced players are more likely to choose remain the same: compute the minimax regret, use its probability distribution to approximate players’ beliefs, and choose the pure strategies that best respond to these beliefs. Table 8.3 presents examples of a game with a single dominant strategy (SD), a prisoner’s dilemma (PD), a pure coordination game (PC), a stag-hunt game (SH) and a symmetric (BS) and an asymmetric (aBS) battle of the sexes.

In accordance with theoretical predictions, the procedure selects the unique Nash equilibrium (NE) in the SD and PD games. For what concerns coordination games the procedure always indicates a single outcome. With this respect it can therefore be considered as a tool for equilibrium selection for games characterized by multiple equilibria. More precisely the procedure selects the Pareto dominant equilibrium in pure coordination games (PC). This is in line with intuition, theory and experimental results. More controversial is the indication for stag hunt games (SH), i.e. games that have a Pareto-dominant NE (more rewarding) and a risk-dominant NE (less risky). The latter is the one indicated by the procedure. For this class of games experimental results provide mixed indications (see, for instance, Hanseny and Selten; 1988; Straub, 1995; Haruvy and Stahl, 2004).

More in general, in games that have at least a NE in pure strategies, if the procedure selects a single outcome, then this outcome is always a NE of the game (SD, PD, PC, SH and aBS). However it may be the case that the procedure does not select any outcome (or better it selects them all), even if pure Nash equilibria exist. This is what happens in the case of asymmetric battle of the sexes (BS) where the expected payoffs of the two strategies conditional on the conjectured beliefs are equal. The situation is different in the asymmetric version of the game (aBS) where the procedure selects the payoff dominant equilibrium. Both predictions are in line with empirical evidences.

Finally notice that the conjectured minimax regret beliefs sometimes happen to be partially incorrect, i.e. they do not confirm in formation in the strategies selected by the opponent (see games SD, BS, aBS). For instance in the SD game the row player expects his opponent to be biased towards playing strategy \( T_B \) but indeed, according to the procedure, player \( B \) chooses strategy \( H_B \). We do not perceive this to be a problem. In fact, we axiomatized the beliefs of inexperienced, unsophisticated and bounded rational players and therefore the possibility that the procedure may assign incorrect beliefs was embedded in our model since the beginning. What matters is that players actually play as if they were best responding to the minimax regret beliefs such that the procedure is effective in predicting players’ behaviour. In the case of the SD game, for instance, player \( A \) chooses his strictly dominant strategy (which is a best response to any possible belief) and player \( B \) best responds with strategy \( H_B \).
8.7 Conclusion

$2 \times 2$ one-shot games remain a fundamental tool for modelling strategic interactions. These games capture the simplest relations (the number of players and strategies is minimal), but still they can be used to describe an uncountable number of situations. No wonder therefore that their study has always attracted a lot of attention. Nevertheless the gap between theoretical models and agents' actual behaviour often happens to be still wide. As a consequence, to predict players' behaviour in one-off interactions remains a problematic issue.

This chapter introduced a simple procedure to be used in forecasting the outcome of $2 \times 2$ one-shot games. Using an axiomatic approach, we looked for a function that may approximate the beliefs of inexperienced and boundedly rational players. First we discussed various proposals connected with the concept of mixed strategy Nash equilibrium and we showed that these functions cannot be expected to mimic players' beliefs satisfactorily. Then we showed that a belief function based on an instrumental use of minimax regret succeeds in this task. The procedure simply allows players to behave as if they were responding to these conjectured minimax regret beliefs. Experimental evidence confirms that the procedure is an effective tool for anticipating the moves of the vast majority of the players.

Notes

1. These papers elicit players' beliefs using a proper quadratic scoring rule, such that for the players 'telling the truth' is optimal.
2. The following game exactly mimics the rod example. And indeed the mixed equilibrium is given by $(1T + 1B)$ for both players.

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3. See, for instance, Osborne and Rubinstein (1994), section 3.2.
6. A notable exception is Goeree and Holt (2004) that presents a model of iterated noisy introspection for one-shot interactions which is then tested over a large number of games.
7. More recently minimax regret has been used to model the behaviour of subjects with bounded rationality (Bergermann and Schlag, 2006) as well as a way to deal with missing data in econometrics (Manski, 2008) and has also appears in the artificial intelligence literature (Brafman and Tennenholtz, 2000).
8. This assumption implies no loss of generality. It is in fact easy to show how the minimax regret proposal satisfy the axioms also in the case of $a < c$ and $d < b$ as well as in the cases with weakly or strictly dominated strategies.
9. These probabilities are analogous to the ones identified by the $H_2$-belief function proposed in Gaillac (2007), although it considers a richer set of axioms but the results of the analysis are analogous.
10. Notice that the procedure considers all the payoffs of the game. Beliefs of player 2 are mimicked by the minimax regret probability distribution of the opponent 1 and thus they depend on the payoffs of the latter. But then, in computing $s_1$' best response, $s_2$'s payoffs are also taken into account.
11. With respect to the original papers strategies will be renamed in order to be consistent with previous sections.
12. The hit rate is described in Verbeek (2004) and used, for instance, in Gneezy and Guth (2003).
13. More precisely $H_2 > T_2$ because $1/320 + 1/40 > 1/40 + 1/80$ and $T_2 > H_2$ because $1/80 + 1/40 + 1/40 + 1/80$.
14. Note that the payoffs structure of the GR games is more complex. Moreover, despite of the fact that games were one-shot, the huge number of strategic situations players had to face makes these data less appropriate to study the behaviour of inexperienced agents.
15. See the graphical example in section 8.4.4; $k = 1$ identifies the symmetric case, i.e. the unique point for which the functions for the minimax regret and for the MSNE intersect.

References


